



## On Mixed Type Duality for Multiobjective Programming Containing Support Functions

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**Abstract.** A mixed type vector dual to a multiobjective programming problem containing support functions is formulated and various duality results are proved under generalized invexity conditions. Special cases are generated from our results.

### 1. Introduction

In [5], Husain *et al.* considered the following multiobjective programming containing support functions

$$(NP) \quad \text{Minimize}(f^1(x) + S(x|C^1), \dots, f^p(x) + S(x|C^p))$$

subject

$$g^j(x) + S(x|D^j) \leq 0, \quad j = 1, 2, \dots, m.$$

Where

- (i)  $f^i : R^n \rightarrow R$  and  $g^j : R^n \rightarrow R$ ,  $j = 1, 2, \dots, m$  are differentiable functions and
- (ii)  $S(\cdot|C^i)$ ,  $i = 1, 2, \dots, p$  and  $S(\cdot|D^j)$ ,  $j = 1, 2, \dots, m$  are support functions of a compact convex set  $C^i$ ,  $i = 1, 2, \dots, p$  and  $D^j$ ,  $j = 1, 2, \dots, m$  in  $R^n$ , to be defined later.

The following Wolfe type dual to the problem (NP) is presented [5]:

$$(WND) \quad \text{Maximize} \left( f^1(u) + u^T z^1 + \sum_{j=1}^m y^j (g^j(u) + u^T w^j), \dots, \right.$$

$$\left. f^p(u) + u^T z^p + \sum_{j=1}^m y^j (g^j(u) + u^T w^j) \right)$$

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subject to

$$\begin{aligned} \sum_{i=1}^p \lambda^i \nabla(f^i(u) + u^T z^i) + \sum_{j=1}^m y^j \nabla(g^j(u) + u^T w^j) &= 0, \\ z^i \in C^i, \quad i = 1, 2, \dots, p, \\ w^j \in D^j, \quad j = 1, 2, \dots, m, \\ y \geq 0, \\ \lambda > 0, \quad \sum_{i=1}^p \lambda^i &= 1. \end{aligned}$$

The problem (WND) is a dual to (NP) assuming that

$$\sum_{i=1}^p \lambda^i (f^i(\cdot) + (\cdot)^T z^i) + \sum_{j=1}^m y^j (g^j(\cdot) + (\cdot)^T w^j)$$

is pseudoinvex with respect to  $\eta$ . The authors in [5] further weakened the invexity required in Wolfe type by constructing the following Mond-Weir type vector dual.

The Mond-Weir vector type dual is the following to (NP):

$$(M-WNP) \quad \text{Maximize } (f^1(u) + u^T z^1, \dots, f^p(u) + u^T z^p)$$

subject to

$$\begin{aligned} \sum_{i=1}^p \lambda^i \nabla(f^i(u) + u^T z^i) + \sum_{j=1}^m y^j \nabla(g^j(u) + u^T w^j) &= 0, \\ z^i \in C^i, \quad i = 1, 2, \dots, p, \\ w^j \in D^j, \quad j = 1, 2, \dots, m, \\ \sum_{j=1}^m y^j (g^j(u) + u^T w^j) &\geq 0, \\ \lambda > 0, y \geq 0. \end{aligned}$$

Husain *et al.* [5] established usual duality theorems under the hypotheses that  $\sum_{i=1}^p \lambda^i \nabla(f^i(\cdot) + (\cdot)^T z^i)$  is pseudoinvex and  $\sum_{j=1}^m y^j \nabla(g^j(\cdot) + (\cdot)^T w^j)$  is quasi-invex with respect to the same  $\eta$ .

In this paper, we propose in the spirit of Husain and Jabeen [4] and Xu [7], a mixed type dual to (NP) to combine the problems (WND) and (M-WNP) and establish various duality theorems under generalized invexity conditions. Special cases are discussed to show that our results extend some earlier results in the literature.

## 2. Pre-requisites

Before stating our multiobjective nonlinear problem, we mention the following conventions for vectors  $x$  and  $y$  in  $n$ -dimensional Euclidian space  $R^n$  to be used

throughout the analysis of this research.

$$\begin{aligned} x < y &\Leftrightarrow x_i < y_i, & i = 1, 2, \dots, n. \\ x \leq y &\Leftrightarrow x_i \leq y_i, & i = 1, 2, \dots, n. \\ x \leq y &\Leftrightarrow x_i \leq y_i, & i = 1, 2, \dots, n, \text{ but } x \neq y \end{aligned}$$

$x \not\leq y$ , is the negation of  $x \leq y$

For  $x, y \in R$ ,  $x \leq y$  and  $x < y$  have the usual meaning.

Before presenting our mixed type dual (Mix D), we mention some definitions of invexity and generalized invexity for easy reference.

**Definition 2.1** (*Invexity*). The function  $\phi : R^n \rightarrow R$  is said to be invex with respect to  $\eta$  at  $\bar{x}$  if there exists a vector function  $\eta(x, \bar{x}) \in R^n$ , such that for all  $x$  and  $\bar{x} \in R^n$

$$\phi(x) - \phi(\bar{x}) \geq \eta(x, \bar{x})^T \nabla \phi(\bar{x}).$$

**Definition 2.2** (*Pseudoinvex*). The function  $\phi : R^n \rightarrow R$  is said to be pseudoinvex with respect to  $\eta$  at  $\bar{x}$  if there exists a vector function  $\eta(x, \bar{x}) \in R^n$ , such that for all  $x$  and  $\bar{x} \in R^n$

$$\eta(x, \bar{x})^T \nabla \phi(\bar{x}) \geq 0$$

implies

$$\phi(x) \geq \phi(\bar{x}).$$

**Definition 2.3** (*Quasi-invex*). The function  $\phi : R^n \rightarrow R$  is said to be quasi-invex with respect to  $\eta$  at  $\bar{x}$  if there exists a vector function  $\eta(x, \bar{x}) \in R^n$ , such that for all  $x$  and  $\bar{x} \in R^n$

$$\phi(x) \leq \phi(\bar{x})$$

implies

$$\eta(x, \bar{x})^T \nabla \phi(\bar{x}) \leq 0.$$

**Definition 2.4** (*Support function*). Let  $K$  be a compact set in  $R^n$ , then the support function of  $K$  is defined by

$$S(x|K) = \max\{x^T v \mid v \in K\}.$$

A support function, being convex everywhere finite, has a subdifferential in the sense of convex analysis. The subdifferential of  $s(x|K)$  is given by

$$\partial S(x|K) = \{z \in K \mid z^T x = S(x|K)\}.$$

For a set  $K$ , the normal cone to  $K$  at a point  $x \in K$  is defines by

$$N_k(x) = \{y \mid y^T(z - x) \leq 0, \text{ for all } z \in K\}.$$

When  $K$  is a compact convex set,  $y$  is in  $N_k(x)$  if and only if  $S(y|K) = x^T y$  i.e.,  $x$  is a subdifferential of  $s$  at  $y$ .

**Definition 2.5** (*Efficient solution*). A feasible solution  $\bar{x}$  is efficient for (NP) if there exist no other feasible  $x$  for (VPE) such that for some  $i \in P = \{1, 2, \dots, p\}$ ,

$$f^i(x) + S(x|C^i) < f^i(\bar{x}) + S(\bar{x}|C^i)$$

and

$$f^j(x) + S(x|C^j) \leq f^j(\bar{x}) + S(\bar{x}|C^j) \quad \text{for all } j \in P, j \neq i.$$

In order to prove the strong duality theorem we will invoke the following lemma due to Chankong and Haimes [1]. In the subsequent analysis we shall denote the set of feasible solutions of the problem (NP) by  $X$ .

**Lemma 2.6.** *A point  $\bar{x} \in X$  is an efficient for (NP), if and only if  $\bar{x} \in X$  solves the following problem:*

$$(P_k(\bar{x})) \quad \text{Minimize } f^k(x) + s(S|C^k)$$

subject to

$$f^i(x) + S(x|C^i) \leq f^i(\bar{x}) + S(\bar{x}|C^i) \quad \forall i \in P$$

$$g^j(x) + S(x|D^j) \leq 0, \quad j = 1, 2, \dots, m.$$

### 3. Mixed type duality

We formulate the following type dual (Mix D) to (NP):

$$(Mix D) \quad \text{Maximize } \left( f^1(u) + u^T z^1 + \sum_{j \in J_\circ} y^j (g^j(u) + u^T w^j), \dots, \right. \\ \left. f^p(u) + u^T z^p + \sum_{j \in J_\circ} y^j (g^j(u) + u^T w^j) \right)$$

subject to

$$(1) \quad \sum_{i=1}^p \lambda^i \nabla(f^i(u) + u^T z^i) + \sum_{j=1}^m y^j \nabla(g^j(u) + u^T w^j) = 0,$$

$$(2) \quad \sum_{j \in J_\alpha} y^j (g^j(u) + u^T w^j) \geq 0, \quad \alpha = 1, 2, \dots, r,$$

$$(3) \quad z^i \in C^i, \quad i = 1, 2, \dots, p,$$

$$(4) \quad w^j \in D^j, \quad j = 1, 2, \dots, m,$$

$$(5) \quad y \geq 0,$$

$$(6) \quad \lambda \in \Lambda,$$

$$\text{where } \Lambda = \left\{ \lambda \in R^p \mid \lambda > 0, \sum_{i=1}^p \lambda^i = 1 \right\}.$$

where  $J_\alpha \subseteq M = \{1, 2, \dots, m\}$ ,  $\alpha = 0, 1, 2, \dots, r$  with  $\bigcup_{\alpha=0}^r J_\alpha = M$  and  $J_\alpha \cap J_\beta = \phi$ , if  $\alpha \neq \beta$ . If  $J_\circ = M$ , then (Mix D) becomes Wolfe type dual considered in [5], if  $J_\circ = \phi$  and  $J_\alpha = M$  for some  $\alpha \in \{1, 2, \dots, r\}$ , then (Mix D) becomes Mond-Weir type dual considered in [5].

**Theorem 3.1** (Weak duality). Let  $\bar{x}$  be feasible for (NP) and  $(u, y, z^1, \dots, z^p, w^1, \dots, w^m, \lambda)$  be feasible for (Mix D). If for all feasible  $(x, u, y, z^1, \dots, z^p, w^1, \dots, w^m, \lambda)$ ,  $\sum_{i=1}^p \lambda^i \nabla(f^i(\cdot) + (\cdot)^T z^i) + \sum_{j \in J_\alpha} y^j (g^j(\cdot) + (\cdot)^T w^j)$  is pseudoinvex and  $\sum_{j \in J_\alpha} y^j (g^j(\cdot) + (\cdot)^T w^j)$ ,  $\alpha = 1, 2, \dots, r$  is quasi-invex with respect to  $\eta$ , then the following cannot hold.

$$(7) \quad f^i(x) + s(x|C^i) \leq f^i(u) + u^T z^i + \sum_{j \in J_\alpha} y^j (g^j(u) + u^T w^j)$$

for all  $i \in \{1, \dots, p\}$ , and

$$(8) \quad f^k(x) + s(x|C^k) < f^k(u) + u^T z^k + \sum_{j \in J_\alpha} y^j (g^j(u) + u^T w^j)$$

for some  $k$ .

**Proof.** Suppose that (7) and (8) hold. Then in view of  $\lambda > 0$  and  $\sum_{i=1}^p \lambda^i = 1$ , (7) and (8) together with  $x^T z^i \leq s(x|C^i)$ ,  $i = 1, 2, \dots, p$  and  $x^T w^j \leq s(x|D^j)$ ,  $j = 1, 2, \dots, m$  and the feasibility for (NP) and (Mix D) imply

$$\begin{aligned} & \sum_{i=1}^p \lambda^i (f^i(x) + x^T z^i) + \sum_{j \in J_\alpha} y^j (g^j(x) + x^T w^j) \\ & < \sum_{i=1}^p \lambda^i (f^i(u) + u^T z^i) + \sum_{j \in J_\alpha} y^j (g^j(u) + u^T w^j) \end{aligned}$$

This in view of the pseudoinvexity of

$$\sum_{i=1}^p \lambda^i (f^i(\cdot) + (\cdot)^T z^i) + \sum_{j \in J_\alpha} y^j (g^j(\cdot) + (\cdot)^T w^j)$$

with respect to  $\eta$ , implies,

$$(9) \quad \eta^T(x, u) \left( \sum_{i=1}^p \lambda^i \nabla(f^i(u) + u^T z^i) + \sum_{j \in J_\alpha} y^j \nabla(g^j(u) + u^T w^j) \right) < 0$$

Since  $\bar{x}$  is feasible for (VP),  $(u, y, z^1, \dots, z^p, w^1, \dots, w^m, \lambda)$  is feasible for (Mix D), and  $x^T w^j \leq s(x|D^j)$ ,  $j = 1, 2, \dots, m$ , we have

$$\sum_{j \in J_\alpha} y^j (g^j(x) + x^T w^j) \leq \sum_{j \in J_\alpha} y^j (g^j(u) + u^T w^j), \quad \alpha = 1, 2, \dots, r.$$

This in view of quasi-invexity of  $\sum_{j \in J_\alpha} y^j (g^j(\cdot) + (\cdot)^T w^j)$ ,  $\alpha = 1, 2, \dots, r$  with respect to  $\eta$ , gives

$$\eta^T(x, u) \left( \sum_{j \in J_\alpha} y^j \nabla(g^j(x) + x^T w^j) \right) \leq 0, \quad \alpha = 1, 2, \dots, r$$

Hence

$$(10) \quad \eta^T(x, u) \nabla \left( \sum_{j \in M-J_\alpha} y^j (g^j(x) + x^T w^j) \right) \leq 0, \quad \alpha = 1, 2, \dots, r$$

Combining (9) and (10), we have

$$(11) \quad \eta^T(x, u) \left( \sum_{i=1}^p \lambda^i \nabla(f^i(u) + u^T z^i) + \sum_{j=1}^m y^j \nabla(g^j(u) + u^T w^j) \right) < 0$$

From the equality constraint of (Mix D), we have

$$(12) \quad \eta^T(x, u) \left( \sum_{i=1}^p \lambda^i \nabla(f^i(u) + u^T z^i) + \sum_{j=1}^m y^j \nabla(g^j(u) + u^T w^j) \right) = 0$$

The relation (12) contradicts (11). Hence the conclusion of the theorem is true.

**Theorem 3.2 (Strong duality).** *Let  $\bar{x}$  be an efficient solution of (NP) and for at least one  $i$ ,  $i \in \{1, 2, \dots, p\}$ ,  $\bar{x}$  satisfies the regularity condition [3] for the problem  $(P_k(\bar{x}))$ . Then there exist  $\lambda \in R^p$  with  $\lambda^T = (\bar{\lambda}^1, \dots, \bar{\lambda}^i, \dots, \bar{\lambda}^p)$ ,  $\bar{y} \in R^m$  with  $\bar{y}^T = (\bar{y}^1, \dots, \bar{y}^i, \dots, \bar{y}^m)$ ,  $z^i \in R^n$ ,  $i = \{1, 2, \dots, p\}$  and  $w^j \in R^n$ ,  $j = 1, 2, \dots, m$  such that  $(x, u, y, z^1, \dots, z^p, w^1, \dots, w^m, \lambda)$  is feasible for (Mix D) and the objectives of (NP) and (Mix D) are equal.*

Further, if the hypotheses of Theorem 1 are satisfied, then

$(x, u, y, z^1, \dots, z^p, w^1, \dots, w^m, \lambda)$  is an efficient solution of (Mix D).

**Proof.** Since  $\bar{x}$  is an efficient solution for  $(P_k(\bar{x}))$ , this implies that there exists  $\xi \in R^p$ ,  $v \in R^m$  with  $\bar{v}^T = (\bar{v}^1, \dots, \bar{v}^i, \dots, \bar{v}^m)$  and  $z^i \in R^n$ ,  $i = \{1, 2, \dots, p\}$  such that

$$(13) \quad \bar{\xi}^k \nabla(f^k(x) + \bar{x}^T \bar{z}^k) + \sum_{\substack{i=1 \\ i \neq k}}^p \bar{\xi}^i \nabla(f^i(x) + \bar{x}^T \bar{z}^i) \\ + \sum_{j=1}^m \bar{y}^j \nabla(g^j(x) + x^T w^j) = 0,$$

$$(14) \quad \sum_{j=1}^m \bar{v}^j \nabla(g^j(x) + x^T w^j) = 0,$$

$$(15) \quad \bar{x}^T \bar{z}^i = S(\bar{x}|C^i), \quad i = 1, 2, \dots, p,$$

$$(16) \quad \bar{x}^T \bar{w}^j = S(\bar{x}|D^j), \quad j = 1, 2, \dots, m,$$

$$(17) \quad z^i \in C^i, \quad i = 1, 2, \dots, p,$$

$$(18) \quad w^j \in D^j, \quad j = 1, 2, \dots, m,$$

$$(19) \quad \xi > 0, \quad \bar{v} \geq 0$$

Dividing (13), (14) and (19) by  $\sum_{i=1}^p \xi^i \neq 0$ , and putting  $\bar{\lambda}^i = \frac{\bar{\xi}^i}{\sum_{i=1}^p \xi^i}$  and  $\bar{y}^i = \frac{\bar{v}^i}{\sum_{i=1}^p \xi^i}$ ,

we have

$$(20) \quad \sum_{i=1}^p \bar{\lambda}^i \nabla(f^i(\bar{x}) + \bar{x}^T \bar{z}^i) + \sum_{j=1}^m \bar{y}^j \nabla(g^j(x) + \bar{x}^T \bar{w}^j) = 0$$

$$(21) \quad \sum_{j=1}^m \bar{y}^j \nabla(g^j(x) + \bar{x}^T \bar{w}^j) = 0$$

$$(22) \quad \lambda > 0, \quad \sum_{i=1}^p \lambda^i = 1, \quad \bar{y} \geq 0$$

The equation (21) implies

$$(23) \quad \sum_{j \in J_0} \bar{y}^j (g^j(x) + \bar{x}^T \bar{w}^j) = 0$$

and

$$(24) \quad \sum_{j \in J_\alpha} \bar{y}^j (g^j(x) + \bar{x}^T \bar{w}^j) = 0, \quad \alpha = 1, 2, \dots, r$$

The relation (20), (22) and (24) imply that  $(x, u, y, z^1, \dots, z^p, w^1, \dots, w^m, \lambda)$  is feasible for (Mix D).

$$f^i(\bar{x}) + \bar{x}^T \bar{z}^i + \sum_{j \in J_0} \bar{y}^j (g^j(x) + \bar{x}^T \bar{w}^j) = f^i(\bar{x}) + S(\bar{x}|C^i), \quad i = 1, 2, \dots, p.$$

This implies the objective of the primal and dual problems are equal.

Further, in view of the assumptions Theorem 1, the efficiency of  $\bar{x}$  for (NP) is immediate.

**Theorem 3.3** (Converse duality). *Let  $(\bar{x}, \bar{y}, \bar{z}^1, \dots, \bar{z}^p, \bar{w}^1, \dots, \bar{w}^m, \bar{\lambda})$  be an efficient solution for (Mix D). Assume that*

(A<sub>1</sub>) *f and g are twice continuously differentiable,*

(A<sub>2</sub>)  *$\nabla f^i(\bar{x}) + \bar{z}^i + \sum_{j \in J_0} \bar{y}^j (\nabla g^j(\bar{x}) + \bar{w}^j)$  are linearly independent,*

(A<sub>3</sub>)  *$\nabla^2(\lambda^T f^i(\bar{x}) + \bar{y}^T g(\bar{x}))$  is positive or negative definite.*

Further, if the assumptions of Theorem 1 are met, then  $\bar{x}$  is an efficient solution.

**Proof.** Since  $(\bar{x}, \bar{y}, \bar{z}^1, \dots, \bar{z}^p, \bar{w}^1, \dots, \bar{w}^m, \bar{\lambda})$  be an efficient solution of (Mix D), then there exist  $\tau \in R^p$ ,  $\beta \in R^n$ ,  $\gamma \in R$  for each  $\gamma$  constraints,  $\eta \in R^p$  with  $\eta^T = (\eta^1, \dots, \eta^i, \dots, \eta^p)$  and  $\mu \in R^m$  such that the following Fritz-John optimality conditions [2] are satisfied,

$$(25) \quad - \sum_{i=1}^p \tau^i \left( \nabla(f^i(\bar{x}) + \bar{x}^T \bar{z}^i) + \sum_{j \in J_0} \bar{y}^j \nabla(g^j(\bar{x}) + \bar{x}^T \bar{w}^j) \right) + \beta^T \nabla^2(\lambda^T f(\bar{x}) + \bar{y}^T g(\bar{x})) - \gamma \sum_{\alpha=1}^r \sum_{j \in J_\alpha} \bar{y}^j \nabla(g^j(x) + \bar{x}^T \bar{w}^j) = 0$$

$$(26) \quad -(\tau^T e)(g^j(\bar{x}) + \bar{x}^T \bar{w}^j) + \beta^T \nabla(g^j(\bar{x}) + \bar{x}^T \bar{w}^j) - \mu^j = 0, \quad j \in J_0$$

$$(27) \quad -\gamma(g^j(\bar{x}) + \bar{x}^T \bar{w}^j + \beta^T \nabla(g^j(\bar{x}) + \bar{x}^T \bar{w}^j)) - \mu^j = 0, \quad j \in J_\alpha, \alpha = 1, \dots, r$$

$$(28) \quad \beta^T \nabla(f(\bar{x}) + \bar{x}^T \bar{z}^i) + \sum_{j \in J_\circ} \bar{y}^j (\nabla g^j(\bar{x}) + \bar{w}^j) - \eta^i = 0$$

$$(29) \quad (\lambda^i \beta - \tau^i \bar{x}) \in N_{C^i}(\bar{z}^i), \quad i = 1, \dots, p$$

$$(30) \quad (\beta - (\tau^T e) \bar{x}) \bar{y}^j \in N_{D^j}(\bar{w}^j), \quad j \in J_\circ$$

$$(31) \quad (\beta - \gamma \bar{x}) \bar{y}^j \in N_{D^j}(\bar{w}^j), \quad j \in J_\alpha, \alpha = i, \dots, r$$

$$(32) \quad \mu^T \bar{y} = 0$$

$$(33) \quad \eta^T \lambda = 0$$

$$(34) \quad \gamma \sum_{j \in J_\alpha} \bar{y}^j \nabla(g^j(\bar{x}) + \bar{x}^T \bar{w}^j) = 0, \quad \alpha = 1, \dots, r$$

$$(35) \quad (\tau, \mu, \eta, \gamma) \geq 0$$

$$(36) \quad (\tau, \beta, \mu, \eta, \gamma) \neq 0$$

Since  $\lambda > 0$ , (33) implies  $\eta = 0$ . Consequently (28) implies

$$(37) \quad \left( \nabla(f^i(\bar{x}) + \bar{x}^T \bar{z}^i) + \sum_{j \in J_\circ} \bar{y}^j (\nabla g^j(\bar{x}) + \bar{w}^j) \right) \beta = 0$$

Using the equality constraint of (Mix D) in (25), we have

$$(38) \quad -\sum_{i=1}^p (\tau^i - \gamma \lambda^i) \left( \nabla f^i(\bar{x}) + \bar{z}^i + \sum_{j \in J_\circ} \bar{y}^j (\nabla g^j(\bar{x}) + \bar{w}^j) \right) + \beta^T \nabla^2(\lambda^T f(\bar{x}) + \bar{y} g(\bar{x})) = 0$$

Postmultiplying (38) by  $\beta$  and then using (37), we have

$$\beta^T \nabla^2(\lambda^T f(\bar{x}) + \bar{y}^T g(\bar{x})) \beta = 0$$

This because of  $(A_3)$ , yields

$$(39) \quad \beta = 0$$

Using (39) along with  $(A_2)$ , we have

$$(40) \quad \tau^i - \gamma \lambda^i = 0, \quad i = 1, 2, \dots, p$$

Suppose  $\gamma = 0$ , then from (40) we have  $\tau = 0$ . Consequently we have from (26) and (27),  $\mu = 0$ .

Thus  $(\tau, \beta, \mu, \eta, \gamma) = 0$ , contradicting (36).

Hence  $\gamma > 0$  and  $\tau > 0$ .

In view of (39), (29), (30) and (31) we have,

$$(41) \quad \bar{x}^T \bar{z}^i = S(\bar{x}|C^i), \quad i = 1, 2, \dots, p$$

$$(42) \quad \bar{x}^T \bar{w}^j = S(\bar{x}|D^j), \quad j = 1, 2, \dots, m$$



From (26) and (27) along with (42) and (35), we have

$$g^j(x) + s(\bar{x}|D^j) \leq 0, \quad j = 1, 2, \dots, m$$

This implies the feasibility of  $\bar{x}$  for (VP).

From (26) and (32), we have

$$\sum_{j \in J_0} \bar{y}^j \nabla(g^j(\bar{x}) + \bar{x}^T \bar{w}^j) = 0$$

In view of this together with (41), we have

$$f^i(\bar{x}) + \bar{x}^T \bar{z}^i + \sum_{j \in J_0} \bar{y}^j (g^j(\bar{x}) + \bar{x}^T \bar{w}^j) = f^i(\bar{x}) + S(\bar{x}|C^i), \quad i = 1, 2, \dots, p$$

This establishes the equality of objective values of (NP).

This in view of the hypothesis of Theorem 1 gives the efficiency of  $\bar{x}$  for (NP).

#### 4. Special cases

In this section, we specialize our problem (NP) and its mixed dual problems (Mix D). As discussed in [6] we may write  $S(x|C^i) = (x^T B^i x)^{\frac{1}{2}}, i = 1, \dots, p$  and  $S(x|D^j) = (x^T E^j x)^{\frac{1}{2}}, j = 1, \dots, m$  and the matrices  $B^i, i = 1, \dots, p$  and  $E^j, j = 1, \dots, m$  are positive semidefinite. Putting these in our problems, we have

$$\begin{aligned} \text{(NP)}_1 \quad & \text{Minimize } (f^1(x) + (x^T B^1 x)^{\frac{1}{2}}, \dots, f^p(x) + (x^T B^p x)^{\frac{1}{2}}) \\ & \text{subject to} \\ & g^j(x) + (x^T E^j x)^{\frac{1}{2}} \leq 0, \quad j = 1, 2, \dots, m \end{aligned}$$

For the dual (Mix D) problem, we get

$$\begin{aligned} \text{(MixD)}_1 \quad & \text{Maximize } \left( f^1(u) + u^T B^1 z^1 + \sum_{j \in J_0} y^j (g^j(u) + u^T E^j w^j) \right) \\ & \left( f^p(u) + u^T B^p z^p + \sum_{j \in J_0} y^j (g^j(u) + u^T E^j w^j) \right) \\ & \text{subject to} \\ & \sum_{i=1}^p \lambda^i \nabla(f^i(u) + u^T B^i z^i) + \sum_{j=1}^m y^j \nabla(g^j(u) + u^T E^j w^j) = 0, \\ & \sum_{j \in J_\alpha} y^j (g^j(u) + u^T E^j w^j) \geq 0, \quad \alpha = 1, 2, \dots, r, \\ & z^T B^i z \leq 1, \quad i = 1, 2, \dots, p, \\ & (w^j)^T E^j w^j \leq 1, \quad j = 1, 2, \dots, m, \\ & \lambda > 0, y \geq 0. \end{aligned}$$

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